

Multiple circle detection based on center-based clustering

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Abstract

The multiple circle detection problem has been considered in the paper on the basis of given data point set $\mathcal{A} \subset \mathbb{R}^2$. It is supposed that all data points from the set \mathcal{A} come from k circles that should be reconstructed or detected. The problem has been solved by the application of center-based clustering of the set \mathcal{A} , i.e. an optimal k -partition is searched for, whose clusters are determined by corresponding circle-centers. Thereby, the algebraic distance from a point to the circle is used. First, an adaptation of the well-known k -means algorithm is given in the paper. Also, the incremental algorithm for searching for an approximate globally optimal k -partition is proposed. The algorithm locates either a globally optimal k -partition or a locally optimal k -partition close to the global one. Since optimal partitions with $2, 3, \dots$ clusters are determined successively in the algorithm, several well-known indexes for determining an appropriate number of clusters in a partition are adopted for this case. Thereby, the Hausdorff distance between two circles is used and adopted. The proposed method and algorithm are illustrated and tested on several numerical examples.

Keywords: Multiple circle detection, Center-based clustering, Globally optimal partition, Approximate optimization, DIRECT

1. Introduction

Clustering or grouping a data set into conceptually meaningful clusters is a well-studied problem in recent literature, and it has practical importance in a wide variety of applications such as medicine, biology, pattern recognition, facility location problem, text classification, information retrieval, earthquake investigation, understanding the Earth's climate, psychology, ranking of municipalities for financial support, business, etc. (Liao et al., 2012; Morales-Esteban et al., 2010; Pintér, 1996; Sabo et al., 2011, 2013; Tan et al., 2006).

A multiple circle detection problem is considered in the paper based on given data set $\mathcal{A} \subset \mathbb{R}^2$. Let us assume that all data from the set $\mathcal{A} = \{a_i = (x_i, y_i) \in \mathbb{R}^2 : i = 1, \dots, m\} \subset \mathbb{R}^2$ come from k circles that should be reconstructed or detected. There are many different approaches for solving this problem in literature, which are most often based on Hough transformation (Ballard, 1981), different heuristic methods (Chung et al., 2012; Cuevas et al., 2012; Kim and Kim, 2001; Qiao and Ong, 2004), RANSAC (Fischler and Bolles, 1981) or fuzzy clustering techniques, that search for the so-called soft or fuzzy partitions (Bezdek et al., 2005; Song et al., 2010; Theodoridis and Koutroumbas, 2009).

A center-based clustering method is applied to solving this problem (Kogan, 2007; Sabo et al., 2013; Tebouille, 2007). The set \mathcal{A} will be grouped into k nonempty disjoint

subsets π_1, \dots, π_k , $1 \leq k \leq m$, such that

$$\bigcup_{i=1}^k \pi_i = \mathcal{A}, \quad \pi_r \cap \pi_s = \emptyset, \quad r \neq s, \quad |\pi_j| \geq 1, \quad j = 1, \dots, k. \quad (1)$$

This partition will be denoted by $\Pi(\mathcal{A}) = \{\pi_1, \dots, \pi_k\}$, and the elements π_1, \dots, π_k of such a partition are called *clusters* in \mathbb{R}^2 .

To each cluster $\pi_j \in \Pi$ a corresponding circle-center $C_j^*(S_j^*, r_j^*)$ with center $S_j^* = (p_j^*, q_j^*)$ and radius r_j^* is associated by solving the following global optimization problem (GOP)

$$(p_j^*, q_j^*, r_j^*) = \operatorname{argmin}_{p, q, r \in \mathbb{R}} \sum_{a_i \in \pi_j} D(C(p, q, r), a_i), \quad (2)$$

where $D(C(p, q, r), a_i)$ represents the distance from the point a_i to the circle C (see Section 2.1).

If the objective function $\mathcal{F}: \mathcal{P}(\mathcal{A}; m, k) \rightarrow \mathbb{R}_+$ is defined on the set of all partitions $\mathcal{P}(\mathcal{A}; m, k)$ of the set \mathcal{A} containing k clusters, then the quality of a partition may be defined and one can search for the *globally optimal k -partition* by solving the following GOP

$$\operatorname{argmin}_{\Pi \in \mathcal{P}(\mathcal{A}; m, k)} \mathcal{F}(\Pi), \quad \mathcal{F}(\Pi) = \sum_{j=1}^k \sum_{a_i \in \pi_j} D(C_j(p_j, q_j, r_j), a_i). \quad (3)$$

Conversely, for a given set of different circles $C_1, \dots, C_k \subset \mathbb{R}^2$, applying the minimal distance principle, the partition $\Pi = \{\pi_1, \dots, \pi_k\}$ of the set \mathcal{A} can be defined in the following way:

$$\pi_j = \{a \in \mathcal{A} : D(C_j, a) \leq D(C_s, a), \quad \forall s = 1, \dots, k, \quad s \neq j\}, \quad (4)$$
$$j = 1, \dots, k,$$

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where one has to pay attention to the fact that every element of the set \mathcal{A} occurs in one and only one cluster. Therefore, the problem of finding an optimal partition of the set \mathcal{A} can be reduced to the following optimization problem

$$\begin{aligned} & \operatorname{argmin}_{C_1, \dots, C_k \subset \mathbb{R}^2} F(C_1, \dots, C_k), \quad (5) \\ F(C_1, \dots, C_k) &= \sum_{i=1}^m \min_{j=1, \dots, k} D(C_j, a_i). \end{aligned}$$

The solution of (3) and (5) coincides (Späth, 1983). In general, the function F can have a great number of independent variables, it does not have to be either convex or differentiable and usually it has several local and global minima. Hence, this becomes a complex GOP.

The adaptation of the well-known k -means algorithm for searching for a locally optimal partition is given in Section 2, where clusters are determined by corresponding circle-centers. In Section 3, a new algorithm for searching for an optimal partition is proposed and an illustrative example is shown. In Section 4, the problem of determining an appropriate number of clusters in a partition is considered. The well-known Davies-Bouldin index, Calinski-Harabasz index, and Simplified Silhouette Width criterion are adopted and used for the case observed in the paper. Finally, in Section 5, several numerical examples are shown.

2. Adaptation of the k -means algorithm

The well-known k -means algorithm (Kogan, 2007; Leisch, 2006; Teboulle, 2007) will be adapted for searching for a locally optimal partition with circles as cluster-centers. The algorithm can be described in two steps which are repeated iteratively:

Algorithm 1. (k closest circles algorithm (KCC))

Step A: For each set of mutually different circles C_1, \dots, C_k the set \mathcal{A} should be divided into k disjoint unempty clusters π_1, \dots, π_k by using the minimal distance principle (4);

Step B: Given a partition $\Pi = \{\pi_1, \dots, \pi_k\}$ of the set \mathcal{A} , one can define the corresponding circle-centers $C_1^*(p_1^*, q_1^*, r_1^*), \dots, C_k^*(p_k^*, q_k^*, r_k^*)$ by solving the following GOPs ($j = 1, \dots, k$)

$$(p_j^*, q_j^*, r_j^*) = \operatorname{argmin}_{p, q, r \in \mathbb{R}} \sum_{a_i \in \pi_j} D(C(p, q, r), a_i). \quad (6)$$

Knowing a good initial approximation, this algorithm can provide an acceptable solution, but in case we do not have a good initial approximation, the KCC-algorithm can be restarted several times with various random initializations (Leisch, 2006).

2.1. The distance from the point to the circle

In order to solve the multiple circle detection problem, it is very important for the distance between a point and a circle to be well defined. This distance will be used by applying the minimal distance principle (4), by determining circle-centers of clusters (6), and by defining the most appropriate number of clusters in a partition in Section 4. Several approaches to determining the distance from the point $a_i = (x_i, y_i) \in \mathbb{R}^2$ to the circle $C(S, r)$ with center $S = (p, q)$ and radius r have been proposed in literature (Ahn et al., 2001; Chernov, 2010; Drezner et al., 2002; Nievergelt, 2002; Song et al., 2010):

$$D_1(C(S, r), a_i) = |||S - a_i|| - r|, \quad (7)$$

$$D_2(C(S, r), a_i) = (||S - a_i|| - r)^2, \quad (8)$$

$$D(C(S, r), a_i) = (||S - a_i||^2 - r^2)^2. \quad (9)$$

The last possibility (9) is called the *algebraic distance*. It is very often used in practical applications (see e.g. Nievergelt (2002); Song et al. (2010)) and for that reason that possibility is also used in our paper.

2.2. Searching for the circle-center of a cluster

The GOPs (6) can have several local and global minima and the corresponding minimizing function is continuous. Note that, if (7) or (8) were used as the distance from the point a_i to the circle C , then the corresponding minimizing function would be a Lipschitz continuous function (see e.g. Sergeyev and Kvasov (2011)). Therefore, for solving GOPs (6) in these cases, some global optimization methods (see e.g. (Jones et al., 1993; Neumaier, 2004; Pintér, 1996; Sergeyev et al., 2001) and corresponding software <http://www.pinterconsulting.com>) can be used. One of the most popular algorithms for solving a GOP for the Lipschitz continuous function is the DIRECT (DIvidingRECTangles) algorithm (Finkel, 2003; Jones et al., 1993; Grbić et al., 2013).

In the case of using the algebraic distance (9), exact solutions of GOPs (6) can be obtained (Chernov, 2010). In the case of using the distance (7) or (8), special methods can also be applied for solving GOPs (6) (Nievergelt, 2002). The random circle consensus technique - RANSAC (Fischler and Bolles, 1981) can also be used for circle fitting.

3. A method for searching for a globally optimal partition

Let $\mathcal{A} \subset [a, b] \times [c, d] \subset \mathbb{R}^2$ be a data points set. A globally optimal k -partition $\Pi^* = \{\pi_1^*, \dots, \pi_k^*\}$ with circle-centers C_1^*, \dots, C_k^* should be determined as a solution of GOP (3), or equivalently (5).

Since objective function (5) is a Lipschitz continuous function, the aforementioned DIRECT algorithm for global optimization can be used. Because of the symmetry property of the function F defined by (5), there are at least $k!$ solutions of this problem. A very efficient special version

of the DIRECT algorithm for symmetric functions is developed in (Grbić et al., 2013). However, complexity of this problem is particularly emphasized if the number of clusters or data points m is large. Instead of solving the GOP (5), motivated by (Bagirov and Ugon, 2005; Bagirov et al., 2011; Scitovski and Scitovski, 2013), an efficient method for searching for a very acceptable approximate of an optimal k -partition is proposed.

First, a sequence of objective functions $F_k: \underbrace{\mathbb{R}^3 \times \dots \times \mathbb{R}^3}_k \rightarrow \mathbb{R}_+$,

$$F_k(C_1, \dots, C_k) = \sum_{i=1}^m \min\{D(C_1, a_i), \dots, D(C_k, a_i)\}, \quad (10)$$

is defined, where $C_j := (S_j, r_j)$, $S_j = (p_j, q_j)$.

For $k = 1$, the problem is reduced to a simple problem of locating a circle on the basis of a given set of points in the plane. The corresponding function $F_1: \mathbb{R}^3 \rightarrow \mathbb{R}_+$ is of the form

$$\begin{aligned} F_1(p_1, q_1, r_1) &= \sum_{i=1}^m D(C_1(p_1, q_1, r_1), a_i) \\ &= \sum_{i=1}^m (\|S_1 - a_i\|^2 - r_1^2)^2, \end{aligned} \quad (11)$$

where $S_1 = (p_1, q_1)$. There exists a number of algorithms and methods for solving this problem (see e.g. (Chernov, 2010; Drezner et al., 2002; Nievergelt, 2002)).

For $k > 1$, an optimal k -partition with circle-centers $C_1^{*(k)}, \dots, C_k^{*(k)}$ is determined by the following incremental algorithm.

Algorithm 2. (Searching for an optimal k -partition)

Step 1: Let $\hat{C}_1, \dots, \hat{C}_{k-1}$ be the solution to the $k-1$ -circle detecting problem and let

$$\begin{aligned} F_{k-1}(\hat{C}_1, \dots, \hat{C}_{k-1}) &= \sum_{i=1}^m \delta_{k-1}^i, \\ \delta_{k-1}^i &= \min\{D(\hat{C}_1, a_i), \dots, D(\hat{C}_{k-1}, a_i)\}, \end{aligned} \quad (12)$$

$$\begin{aligned} \Phi_k(p, q, r) &:= F_k(\hat{C}_1, \dots, \hat{C}_{k-1}, C(p, q, r)) = \\ &= \sum_{i=1}^m \min\{\delta_{k-1}^i, D(C(p, q, r), a_i)\}. \end{aligned} \quad (13)$$

Step 2: By using the DIRECT algorithm for global optimization determine

$$\hat{C}_k \in \underset{\substack{(p,q) \in [a,b] \times [c,d] \\ r \in [0, r_{max}]}}{\operatorname{argmin}} \Phi_k(p, q, r), \quad (14)$$

where r_{max} is the maximum radius of circle C , which can be expected in rectangle $[a, b] \times [c, d]$.

Step 3: By using the KCC-algorithm with initial approximations $\hat{C}_1, \dots, \hat{C}_k$ determine new centers $C_1^{*(k)}, \dots, C_k^{*(k)}$.

An important assumption of the algorithm is usage of an efficient algorithm for solving a GOP in Step 2. The DIRECT algorithm for global optimization of a Lipschitz continuous function was proposed by (Jones et al., 1993). The algorithm operates by systematically dividing the box domain of the objective function into hyperrectangles, and evaluating its values in their centers. There are two phases to an iteration of DIRECT; first, hyperrectangles are identified as potentially optimal, i.e., they have potential to contain a global solution. The second phase of an iteration is to divide potentially optimal hyperrectangles into smaller hyperrectangles. The objective function is evaluated in the centers of new hyperrectangles.

By increasing the number of circle-centers (Step 2) and by applying the KCC algorithm (Step 3), the objective function value does not increase. Therefore,

$$\begin{aligned} F_1(C_1^{*(1)}) &\geq F_2(C_1^{*(1)}, \hat{C}_2) \geq F_2(C_1^{*(2)}, C_2^{*(2)}) \geq \dots \\ &\geq F_k(C_1^{*(k-1)}, \dots, \hat{C}_k) \geq F_k(C_1^{*(k)}, \dots, C_k^{*(k)}) = F_k^*, \end{aligned}$$

and, according to (Bagirov and Ugon, 2005), the maximum number of clusters k_{max} that makes sense to be calculated using Algorithm 2 is determined by

$$\frac{F_{k_{max}-1}^* - F_{k_{max}}^*}{F_1^*} < \epsilon, \quad (15)$$

for some small $\epsilon > 0$. Namely, in that case the relative reduction of the objective function value for $k \geq k_{max}$ is less than ϵ .

Remark 1. Note that the partitions obtained in this way can unfortunately not be said to be globally optimal, but numerous calculations in Section 5 show that the partitions obtained in this way are either globally optimal or locally optimal partitions close to globally optimal ones and therefore acceptable in practical applications. In what follows, the partition obtained by Algorithm 2 will simply be called an optimal partition.

Note also that by using Algorithm 2 an optimal partition for each $k \leq k_{max}$ is obtained, which makes it possible to decide on the appropriate number of clusters in a partition by adaptation of various well-known indexes (see Section 4).

In the following example the proposed Algorithm 2 will be illustrated on synthetic data.

Example 1. The set of 5 circles $C = \{C_i = S_i + r_i(\cos t, \sin t) : t \in [0, 2\pi], i = 1, \dots, 5\}$ is given in the plane. In the neighborhood of the i -th circle $n_i \sim \mathcal{U}(30, 50)$ random points are generated by using binormal random additive errors with mean vector $\mathbf{0} \in \mathbb{R}^2$ and the covariance matrix $.05\mathbf{I}$, where $\mathbf{I} \in \mathbb{R}^{2 \times 2}$ is the identity matrix. In this way, a data point set $\mathcal{A} = \{a_i \in \mathbb{R}^2 : i = 1, \dots, m\}$, with $m = 47 + 34 + 43 + 34 + 43$ random points is obtained.

The algorithm starts with initial circle $\hat{C}_1 = ((3, 3), 1)$ (see Fig. 1a left). By using the DIRECT algorithm the circle \hat{C}_2 is determined in Step 2 (see also Fig. 1a left) and

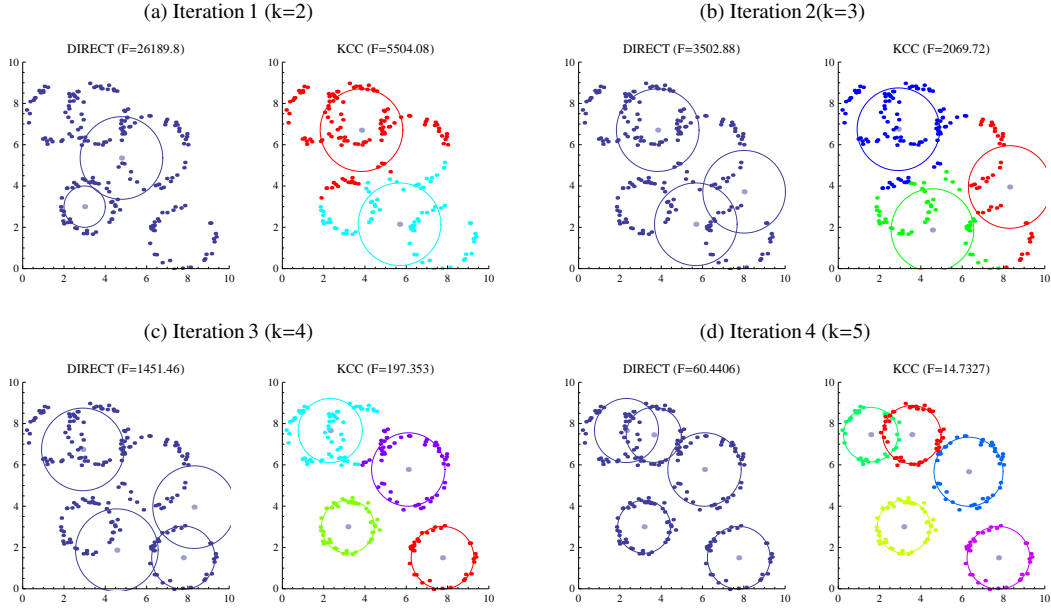


Figure 1: The flow of the iterative process of Algorithm 2

after that, in *Step 3*, circles $C_1^{*(2)}, C_2^{*(2)}$ are obtained by using the KCC algorithm (see Fig. 1a right). A further flow of the iterative process provided by Algorithm 2 is illustrated in Fig. 1b-d. In the end of the iterative process the set of recognized circles $C^* = \{C_i^{*(5)}; i = 1, \dots, 5\}$ is obtained. Objective function values during the iterative process of Algorithm 2 are also shown in Fig. 1.

The quality of recognition can be measured by the Hausdorff distance between sets C and C^*

$$d_H(C, C^*) = \max\{\max_{C_r \in C} \min_{C_s^* \in C^*} \delta(C_r, C_s^*), \max_{C_s^* \in C^*} \min_{C_r \in C} \delta(C_r, C_s^*)\}. \quad (16)$$

In this example, $d_H(C, C^*) = 0.05465$.

4. The most appropriate number of clusters

Determining the number of clusters automatically has been one of the most difficult problems in data clustering processes (Gan et al., 2007; Kaufman and Rousseeuw, 2005; Kogan, 2007; Vendramin et al., 2009). In simple cases, the number of clusters in a partition is determined by the nature of the problem itself. If the number of clusters in a partition is not given in advance, then it is natural to search for an optimal partition which consists of clusters that are as compact and relatively strongly separated as possible. For determining the most appropriate number of clusters in a partition it is possible to find numerous indexes that point to that number. Some of these indexes are particularly analyzed and adapted in this paper. For this purpose, in addition to the distance measure between a point and a circle, it is important to have a good measure for the distance between two circles.

4.1. The distance between two circles

The distance between two circles $C_1 = (S_1, r_1), C_2 = (S_2, r_2)$ in the plane can be defined as the Hausdorff dis-

tance (Chen et al., 2010; Jüttler, 2001)

$$d_H(C_1, C_2) = \max\{\max_{Y \in C_2} \min_{X \in C_1} \|X - Y\|, \max_{X \in C_1} \min_{Y \in C_2} \|X - Y\|\}, \quad (17)$$

where $\|\cdot\|$ denotes the Euclidean distance.

Proposition 1. For any two circles $C_1 = (S_1, r_1), C_2 = (S_2, r_2)$ in the plane the Hausdorff distance is given by

$$d_H(C_1, C_2) = \|S_1 - S_2\| + |r_2 - r_1| =: \delta(C_1, C_2). \quad (18)$$

4.2. Adaptation of some of the known indexes

Let $\mathcal{A} \subset \mathbb{R}^2$ be the set and $\Pi^* = \{\pi_1^*, \dots, \pi_k^*\}$ an optimal k -partition with circle-centers $C_1^*, \dots, C_k^* \in \mathbb{R}^2$. Thereby $C_j^* = (S_j^*, r_j^*)$, where $S_j^* = (p_j^*, q_j^*)$ is the center and r_j^* the radius of the circle C_j^* .

(a) *Davies - Bouldin index* in the standard case for the optimal partition $\Pi^* = \{\pi_1^*, \dots, \pi_k^*\}$ of the set \mathcal{A} with k clusters π_1^*, \dots, π_k^* and the corresponding centers $c_1^*, \dots, c_k^* \in \mathbb{R}^2$ is defined by (see e.g. (Vendramin et al., 2009))

$$DB(k) = \frac{1}{k} \sum_{j=1}^k \max_{\substack{s=1, \dots, k \\ s \neq j}} \frac{V(\pi_j^*) + V(\pi_s^*)}{\|c_j^* - c_s^*\|}, \quad (19)$$

where

$$V(\pi_j^*) = \frac{1}{|\pi_j^*|} \sum_{a_s \in \pi_j^*} \|c_j^* - a_s\|.$$

That is why the *Davies - Bouldin index for circle-centers* will be defined analogously to (19) by

$$DBC(k) = \frac{1}{k} \sum_{j=1}^k \max_{\substack{s=1, \dots, k \\ s \neq j}} \frac{\hat{V}(\pi_j^*) + \hat{V}(\pi_s^*)}{\delta^2(C_j^*, C_s^*)}, \quad (20)$$

where

$$\hat{V}(\pi_j^*) = \frac{1}{|\pi_j^*|} \sum_{a_s \in \pi_j^*} \sqrt{D(C_j^*, a_s)},$$

and $\delta(C_j^*, C_s^*)$ is the distance between circles C_j^*, C_s^* defined by (18) in *Proposition 1*. Let us notice that in the definition of the index (20) the property of non-dimensionality has been retained, since nominators and denominators have the same dimension (of squared length).

More compact and better separated clusters in an optimal partition will result in a lower DBC index.

(b) *Calinski-Harabasz index* in the standard case for the optimal partition $\Pi^* = \{\pi_1^*, \dots, \pi_k^*\}$ of the set A with k clusters π_1^*, \dots, π_k^* and the corresponding centers $c_1^*, \dots, c_k^* \in \mathbb{R}^2$ is defined by (see e.g. (Vendramin et al., 2009))

$$\text{CH}(k) = \frac{\mathcal{G}(c_1^*, \dots, c_k^*)/(k-1)}{\mathcal{F}(c_1^*, \dots, c_k^*)/(m-k)}, \quad (21)$$

where

$$\mathcal{F}(c_1^*, \dots, c_k^*) = \sum_{j=1}^k \sum_{a_i \in \pi_j^*} \|c_j^* - a_i\|^2,$$

$$\mathcal{G}(c_1^*, \dots, c_k^*) = \sum_{j=1}^k |\pi_j^*| \|c_j^* - c^*\|^2,$$

where $c_j^* = (1/|\pi_j|) \cdot \sum_{a_i \in \pi_j} a_i$ are centroids of clusters π_j , and $c^* = (1/|\mathcal{A}|) \cdot \sum_{a_i \in \mathcal{A}} a_i$ is a centroid of the entire set \mathcal{A} .

That is why the *Calinski-Harabasz index for circle-centers* will be defined analogously to (21) by

$$\text{CHC}(k) = \frac{\hat{\mathcal{G}}(C_1^*, \dots, C_k^*)/(k-1)}{\hat{\mathcal{F}}(C_1^*, \dots, C_k^*)/(m-k)}, \quad (22)$$

where

$$\begin{aligned} \hat{\mathcal{F}}(C_1^*, \dots, C_k^*) &= \sum_{j=1}^k \sum_{a_i \in \pi_j^*} \sqrt{D(C_j^*(S_j^*, r_j^*), a_i)}, \\ \hat{\mathcal{G}}(C_1^*, \dots, C_k^*) &= \sum_{j=1}^k |\pi_j^*| \delta^2(C^*(S^*, r^*), C_j^*(S_j^*, r_j^*)) = \\ &= \sum_{j=1}^k |\pi_j^*| (\|S^* - S_j^*\| + |r^* - r_j^*|)^2 \end{aligned}$$

where $C_j^* = (S_j^*, r_j^*)$ are center-circles of clusters π_j^* , and $C^* = (S^*, r^*)$ is a center-circle of the entire set \mathcal{A} . In that way, by definition (22), the property of non-dimensionality for CHC index has also been retained.

More compact and better separated clusters in an optimal partition will result in a greater CHC index.

(c) *Simplified Silhouette Width Criterion* in the standard case for the optimal partition $\Pi^* = \{\pi_1^*, \dots, \pi_k^*\}$ of the set \mathcal{A} with k clusters π_1^*, \dots, π_k^* and the corresponding centers $c_1^*, \dots, c_k^* \in \mathbb{R}^2$ is defined in the following way (see

e.g. (Kaufman and Rousseeuw, 2005; Vendramin et al., 2009)). For each $a_i \in \mathcal{A} \cap \pi_r^*$ the numbers

$$\alpha_{ir} = \|c_r^* - a_i\|, \quad \beta_{ir} = \min_{\substack{s=1, \dots, k \\ s \neq r}} \|c_s^* - a_i\|,$$

$$s_i = \frac{\beta_{ir} - \alpha_{ir}}{\max\{\alpha_{ir}, \beta_{ir}\}},$$

are calculated and the *Simplified Silhouette Width Criterion* is defined as the average of s_i : $\text{SSW}(k) = (1/|\mathcal{A}|) \cdot \sum_{a_i \in \mathcal{A}} s_i$.

Analogously, the *Simplified Silhouette Width criterion for circle-centers* will be defined. For each $a_i \in \mathcal{A} \cap \pi_r^*$ the numbers

$$\hat{\alpha}_{ir} = D(C_r^*, a_i), \quad \hat{\beta}_{ir} = \min_{\substack{s=1, \dots, k \\ s \neq r}} D(C_s^*, a_i),$$

$$\hat{s}_i = \frac{\hat{\beta}_{ir} - \hat{\alpha}_{ir}}{\max\{\hat{\alpha}_{ir}, \hat{\beta}_{ir}\}},$$

are calculated and the *Simplified Silhouette Width Criterion for circle-centers* is defined as the average of \hat{s}_i : $\text{SSWC}(k) = (1/|\mathcal{A}|) \cdot \sum_{a_i \in \mathcal{A}} \hat{s}_i$.

More compact and better separated clusters in an optimal partition will result in a greater SSWC number.

5. Numerical examples

Example 2. *The iterative process of Algorithm 2 for $k \leq 8$ has been implemented for the data from Example 1.*

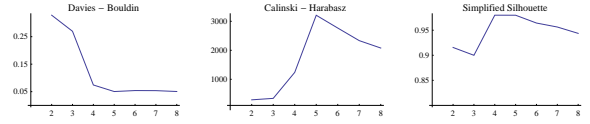


Figure 2: Indexes for the data from Example 2

The configuration of the obtained circles in iterations after Step 2 (application of the DIRECT algorithm) and Step 3 (application of the KCC algorithm) for the first four iterations is shown previously in Fig 1. The graphs of the used indexes are shown in Fig. 2. The DBC index gives a certain priority to the partition with 5 clusters, the CHC index gives a clear priority to the same partition, whereas the SSWC criterion gives an insignificant priority to the partition with 4 clusters compared to the partition with 5 clusters.

Therefore, it can be said that all the indexes used in this example have pointed to the partition with 5 clusters as the most appropriate partition, whereby the Hausdorff distance between original and recognized circles is sufficiently small ($d_H(C, C^*) = 0.05$).

Example 3. *Set C of 4 concentric circles with common center $S = (5, 5)$ and radii $r_i = .75, 1.5, 2.5, 3.5$ is given in the plane. Data point set \mathcal{A} is constructed similarly to Example 1 and it consists of $m = 40 + 40 + 43 + 43$ random points around circles.*

(a) Well separated circles (b) Mutually intersecting circles

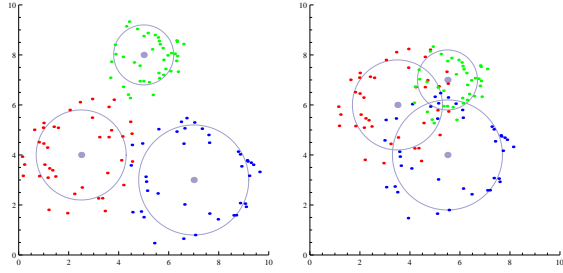
Figure 5: Data points from Example 4 with $\mu = 40$ and $\sigma = 0.4$

Table 1: Percentage of recognition of Algorithm 2 (basic numbers) and the Hough Transform (numbers in parentheses)

	Well separated circles			Mutually intersecting circles		
	$\sigma = 0.1$	$\sigma = 0.2$	$\sigma = 0.4$	$\sigma = 0.1$	$\sigma = 0.2$	$\sigma = 0.4$
$\mu = 20$	85 (1)	72 (1)	52 (0)	68 (5)	52 (3)	41 (2)
$\mu = 40$	97 (4)	94 (4)	64 (2)	89 (3)	84 (5)	60 (2)
$\mu = 60$	99 (1)	96 (1)	86 (2)	87 (4)	86 (2)	77 (2)

in Table 1 show the number of experiments in which Algorithm 2 recognized the set of circles ($d_H(C, C^*) < 0.1$) well. The number of such cases is large and it depends on the number of data μ along the circle and parameter σ in the covariance matrix. This proves a very good performance of the proposed algorithm. The numbers in parentheses show the number of experiments in which the Hough Transform for circle detection recognized the circles well. The number of such cases is very small. It turned out that the Hough Transform is not designed for such case of a relatively small number of data points. Only in special cases, when it was possible to adjust certain values of a Matlab option of “imfindcircles” (e.g., sensitivity factor), a good reconstruction of circles was obtained.

A similar experiment is repeated for three mutually intersecting circles (see Fig. 5b), and the obtained results are similar.

6. Conclusions

The method and the algorithm for solving the multiple circle detection problem on the basis of the given data point set $\mathcal{A} \subset \mathbb{R}^2$ gives either a globally optimal k -partition or a locally optimal k -partition close to the global one. It can be seen that center-circles obtained in this way present a good approximation of the original circles. The method has been tested on several numerical examples with synthetic data on the basis of our own software done by *Mathematica*. The Hausdorff distance between original and recognized circles shows an acceptable quality of reconstruction. It is important to note that minimizing functions (12)–(13) can have several local and global minima and that, precisely for this reason, it is necessary to use some global optimization method. Since minimizing

functions are Lipschitz continuous functions, the DIRECT algorithm was found to be very useful.

Assuming that we have a good approximation of the optimal partition, indexes adapted to the multiple circle detection problem recognized the appropriate number of clusters in a partition well. Thereby, algebraic distance (9) from a point to a circle and the Hausdorff distance between two circles have been used in the paper. If any other distance measure from a point to a circle and the distance between two circles were used, then the corresponding indexes would have to be redefined. Also, some other well-known indexes could be considered.

Finally, the considered problem can be generalized such that ellipses in the plane $E_j = \{x \in \mathbb{R}^2: \|x - C_j\|_{A_j} = 1\}$ are considered instead of circles, where C_j is the center of ellipse E_j and $A_j > 0$ is a positive definite symmetric matrix which determines the major and the minor axes lengths as well as the orientation of the ellipse. In that way, the ellipse is determined by its 5 parameters. The problem can be further generalized to a hyperellipse in \mathbb{R}^n (Bezdek et al., 2005).

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Appendix

Proof of Proposition 1

Without loss of generality, it can be supposed that

$$S_1 = (0, 0), \quad S_2 = (d, 0), \quad d \geq 0, \quad r_2 \geq r_1 > 0. \quad (23)$$

It should be proved that

$$d_H(C_1, C_2) = d + r_2 - r_1. \quad (24)$$

Let us notice the following points on circle C_1 , i.e. C_2

$$\begin{aligned} X_1 &= (-r_1, 0), \quad X_2 = (r_1, 0) \in C_1, \\ Y_1 &= (d - r_2, 0), \quad Y_2 = (d + r_2, 0) \in C_2. \end{aligned}$$

The Hausdorff distance between circles C_1, C_2 is attained precisely on one of the pairs of points (X_i, Y_j) , $i, j = 1, 2$. That is why (17) becomes

$$\begin{aligned} d_H(C_1, C_2) &= \max\{ \\ &\max\{\min\{\|X_1 - Y_1\|, \|X_2 - Y_1\|\}, \min\{\|X_1 - Y_2\|, \|X_2 - Y_2\|\}\}, \\ &\max\{\min\{\|X_1 - Y_1\|, \|X_1 - Y_2\|\}, \min\{\|X_2 - Y_1\|, \|X_2 - Y_2\|\}\} \}. \end{aligned} \quad (25)$$

In particular, if $S_1 = S_2$, i.e. $d = 0$, the Proposition

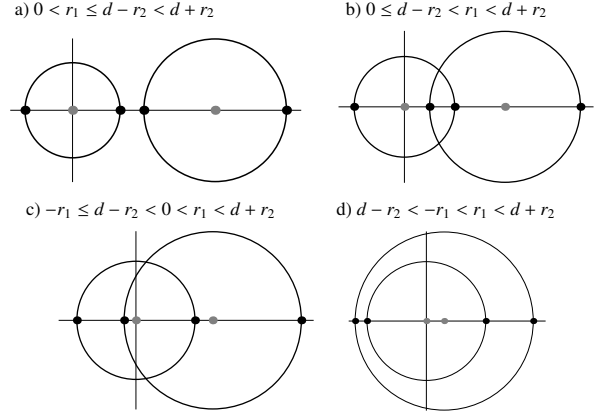


Figure 6: Possible positions of circles C_1, C_2 with conditions in (23)

assertion is true, because then $d_H = r_2 - r_1$.

Therefore, suppose that $d > 0$ and consider all possible positions of circles C_1, C_2 , by which conditions (23) remain fulfilled (see Fig. 6):

- $0 < r_1 \leq d - r_2 < d + r_2$
- $0 \leq d - r_2 < r_1 < d + r_2$
- $-r_1 \leq d - r_2 < 0 < r_1 < d + r_2$
- $d - r_2 < -r_1 < r_1 < d + r_2$

It should be proved that equality (24) holds in each of the mentioned cases. Only the proof of assertion in case *a*) will be considered. In other cases the proof is analogue.

Suppose that $0 < r_1 \leq d - r_2 < d + r_2$ (see Fig. 6a). Then there holds

$$\begin{aligned} \min\{\|X_1 - Y_1\|, \|X_2 - Y_1\|\} &= \|X_2 - Y_1\| = d - r_2 - r_1, \\ \min\{\|X_1 - Y_2\|, \|X_2 - Y_2\|\} &= \|X_2 - Y_2\| = d + r_2 - r_1, \end{aligned}$$

from which it follows that

$$\begin{aligned} \max\{\min\{\|X_1 - Y_1\|, \|X_2 - Y_1\|\}, \min\{\|X_1 - Y_2\|, \|X_2 - Y_2\|\}\} &= \\ &= d + r_2 - r_1. \end{aligned} \quad (26)$$

Similarly, from

$$\begin{aligned} \min\{\|X_1 - Y_1\|, \|X_1 - Y_2\|\} &= \|X_1 - Y_1\| = d - r_2 + r_1, \\ \min\{\|X_2 - Y_1\|, \|X_2 - Y_2\|\} &= \|X_2 - Y_1\| = d - r_2 - r_1, \end{aligned}$$

follows

$$\begin{aligned} \max\{\min\{\|X_1 - Y_1\|, \|X_1 - Y_2\|\}, \min\{\|X_2 - Y_1\|, \|X_2 - Y_2\|\}\} &= \\ &= d - r_2 + r_1. \end{aligned} \quad (27)$$

By using (26) and (27) in (25) the required equality (24) for case *a*) is obtained.